



Mind the Gap

reflections on benchmarking in the global era

Ofer M. Shir (Tel-Hai College & Migal Institute, ISRAEL)
ofersh@telhai.ac.il

Dagstuhl Seminar 23251, June 2023



SCHLOSS DAGSTUHL
Leibniz-Zentrum für Informatik

“Challenges in Benchmarking Optimization Heuristics”

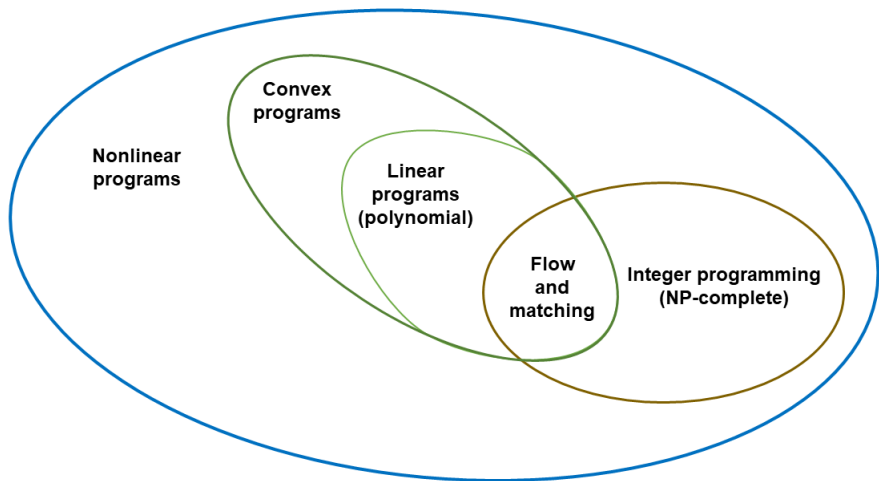
Consensus on Benchmarking Compilation?

The **human factor plays a crucial role in such a process**. Formulation of a test-suite may involve three types of scholars: theoreticians, algorithms' designers, and practitioners:

- (i) **theoreticians** naturally favor analyzable functions
- (ii) **algorithms' engineers** may prefer families of functions that are successfully treated by their designs
- (iii) **practitioners** may have the best insights into which functions most accurately represent real-world problems (thus having their biased preferences)

A proper balance should be made amongst those three parties to effectively compile a test-suite meaningful to a broad audience.¹

Yet Another Perspective: Operations Research as the Optimization Complement²



Opportunities: OR Gaps

Could we perhaps make a difference by actually addressing the weak spots of OR as our benchmarks?

Algorithms-wise, hybrids do make a difference while outperforming white-box solvers – see, e.g., MATHEURISTICS and Blum+Raidl.³

Existing opportunity:

Aim for compiling an “OpenEC” platform, with a unified interface, for optimizing test-functions that are beyond the reach of OR.

Continuous Opportunity (SO): The Quantum Sphere⁴

Given a control phase $\phi(\Omega)$ and a dispersion term $k(\Omega)$, the spectral intensity $I_2(\Omega)$ is formulated by means of the *spectral* terms:

$$\begin{aligned} E_1(\Omega) &= A(\Omega) \exp i [\phi(\Omega) + k(\Omega)] \\ E_2(\Omega) &= E_1(\Omega) * E_1(\Omega) = \int_{-\infty}^{\infty} E_1(\Omega') \cdot E_1(\Omega - \Omega') d\Omega' \\ I_2(\Omega) &= |E_2(\Omega)|^2. \end{aligned} \quad (1)$$

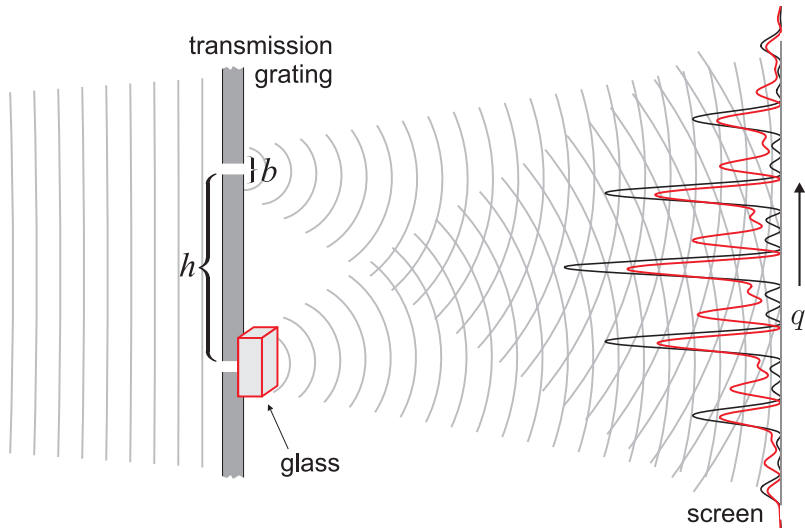
The standard Quantum Sphere is defined as follows:

$$\boxed{QS(\phi, k) = \int_{-\infty}^{\infty} I_2(\Omega) d\Omega \longrightarrow \max}. \quad (2)$$

When $k = 0$, QS is maximized by any phase function linear in the frequency Ω : $\operatorname{argmax}_{\phi(\Omega)} \{QS(\phi, k = 0)\} \equiv a\Omega + b \quad \forall a, b$ (and in particular by a constant phase, $a = 0, \forall b$).

\implies The decision variables ϕ are meant to compensate for the function $k(\Omega)$ over the **periodic domain** $[0, 2\pi]^D$.

Continuous Opportunity (MO): Diffraction Grating⁵



DG Formulation

Introducing a basic set of optical test-problems for Pareto optimization, scalable in dimension and subject to a collection of defining parameters for setting the Pareto front's curvature.

Given an optical setup of n slits, and given phases $\vec{\varphi} \in [0, 2\pi]^n$, the intensity on a screen point positioned at q reads:

$$I_{DG}(q, \vec{\varphi}) = \frac{1}{n^2} \text{sinc}^2\left(\frac{qb}{2}\right) \cdot \left| \sum_{k=0}^{n-1} \exp(iqhk) \cdot \exp(i\varphi_k) \right|^2 \quad (3)$$

Given a set of m competing points on the screen, described by a *position vector* $\vec{q} \in \mathbb{R}^m$, the m -objective DG problem becomes:

$$\vec{f}(\vec{q}, \vec{\varphi}) = \begin{pmatrix} I_{DG}(q_1, \vec{\varphi}) \\ I_{DG}(q_2, \vec{\varphi}) \\ \vdots \\ I_{DG}(q_m, \vec{\varphi}) \end{pmatrix} \longrightarrow \max \quad (4)$$

Discrete Opportunity: Nonlinear Nonconvex Objective⁶

The Low-Autocorrelation Binary Sequence (LABS) problem is a hard combinatorial optimization problem with practical applications in electrical engineering. It is incorporated within the IOHprofiler.

Given a sequence of length n , $S := (s_1, \dots, s_n)$ with $s_i = \pm 1$,

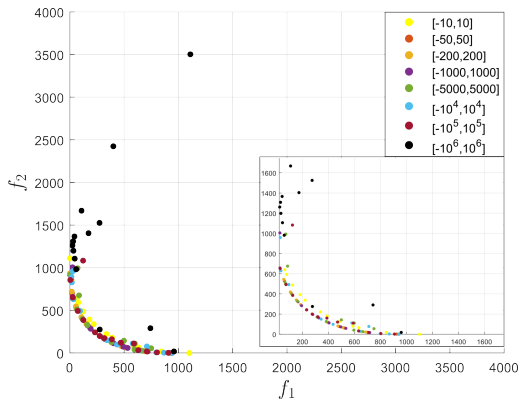
$$\text{[LABS] maximize } \frac{n^2}{2E(S)}$$

subject to:

$$E(S) := \sum_{k=1}^{n-1} \left(\sum_{i=1}^{n-k} s_i \cdot s_{i+k} \right)^2$$
$$s_i \in \{-1, +1\} \quad \forall i \in \{1 \dots n\}$$

(5)

Mixed-Integer Opportunity: White-Box Solvers' Deficiency on moMIQP



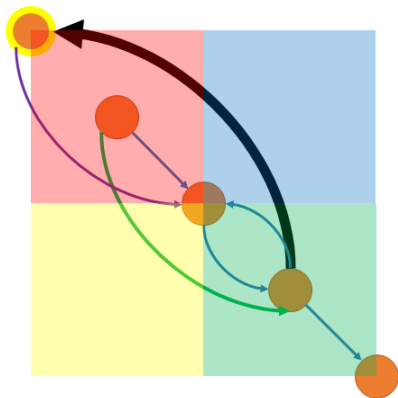
Shir and Emmerich, “Comparing White-Box to Black-Box Solvers over Multi-Objective Mixed-Integer Quadratic Models”, *submitted* (2023).

Complexity of Unbounded Problems

Theorem [Jeroslow, 1973]⁷

A quadratically-constrained mixed-integer unbounded problem is undecidable.

- ① A generalized unbounded MIQP **cannot be linearized** (even if the integers $x_i : i \in I$ may be linearized using auxiliary binaries, the multiplication of two (loosely bounded) decision variables within \vec{x} cannot be precisely linearized (rather could be piecewise approximated upon separation)).
- ② Despite this theoretical complexity, there has been **much practical progress in treating Quadratically-Constrained problems**⁸ (e.g., by reformulation to a bilinear programming problem with integer variables, or by diverting to Mixed-Integer Second-Order Cone Programming when the model permits).



Discussion

danke

Bibliography

- [1] O. M. Shir, C. Doerr, and T. Bäck, “Compiling a benchmarking test-suite for combinatorial black-box optimization: a position paper,” in *Proc. of Genetic and Evolutionary Computation Conference (GECCO'18), Companion*, pp. 1753–1760, ACM, 2018.
- [2] O. M. Shir and T. Bäck, “Sequential experimentation by evolutionary algorithms,” in *Proceedings of the Genetic and Evolutionary Computation Conference Companion, GECCO '22*, (New York, NY, USA), p. 1450–1468, Association for Computing Machinery, 2022.
- [3] C. Blum and G. R. Raidl, *Hybrid Metaheuristics: Powerful Tools for Optimization*. Artificial Intelligence: Foundations, Theory, and Algorithms, Switzerland: Springer International Publishing, 2016.
- [4] O. M. Shir, X. Xing, and H. Rabitz, “Multi-level evolution strategies for high-resolution black-box control,” *Heuristics*, vol. 27, p. 1021–1055, 2021.
- [5] O. M. Shir, J. Roslund, Z. Leghtas, and H. Rabitz, “Quantum control experiments as a testbed for evolutionary multi-objective algorithms,” *Genetic Programming and Evolvable Machines*, pp. 1–47, 2012.
- [6] C. Doerr, F. Ye, N. Horesh, H. Wang, O. M. Shir, and T. Bäck, “Benchmarking discrete optimization heuristics with IOHprofiler,” *Applied Soft Computing*, vol. 88, p. 106027, 2020.
- [7] R. G. Jeroslow, “There cannot be any algorithm for integer programming with quadratic constraints,” *Operations Research*, vol. 21, no. 1, pp. 221–224, 1973.
- [8] C. Blik, P. Bonami, and A. Lodi, “Solving Mixed-Integer Quadratic Programming problems with IBM-CPLEX: a progress report,” in *Proceedings of the 26th RAMP Symposium, Hosei University, Tokyo*, 2014.