

Compiling a Benchmarking Test-Suite for Combinatorial Black-Box Optimization: A Position Paper

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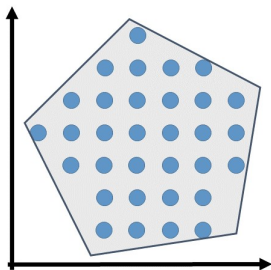
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Benchmarking Discrete Optimization: Fundamentals



Introduction

- Combinatorial Optimization (CO),

$$\mathcal{P} := \left(\mathcal{S}, f : \mathcal{S} \rightarrow \mathbb{R}^+ \right),$$

is defined by a finite set \mathcal{S} with an objective function f assigning a non-negative value to any of its elements $s \in \mathcal{S}$.

[**Seminal work:** “Combinatorial Optimization” by Papadimitriou & Steiglitz]

- Problem-solving in practice: Mathematical Programming (MP) / Operations Research, *versus* Randomized Search Heuristics (RSHs) / Soft Computing
- There are no common grounds for performance comparisons between RSHs to MP solvers when targeting similar CO problems.
- **Benchmarking RSHs on CO-problems is an open issue.**

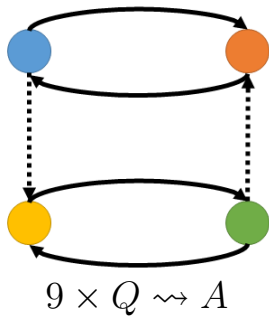
The Role of Benchmarking

- How does the algorithm perform on different **classes of problems** and how does its performance compare to that of other approaches?
- Which **problem features** possess the strongest impact on the accuracy and/or the convergence speed, and how this dependency may be quantified? E.g., *modality* of a problem, its *separability*, the degree of *constraints*, and its *monotonicity*.
- How does the performance **scale with increasing problem complexity** (i.e., dimensionality, cardinality of categories per a decision variable, etc.)?
- How **sensitive** is a given algorithm with respect to small changes in the problem instance or the algorithmic components?

Starting Point

Current focus: formulating a set of benchmark problems and/or a test-suite for CO problems when treated as black-boxes by RSHs.

- Previously proposed guidelines for black-box benchmarking (Whitley et al., 1996):
 - (A) *“Test suites should contain problems that are resistant to hill-climbers”.*
 - (B) *“Test suites should contain problems that are non-linear, non-separable, and non-symmetric”.*
 - (C) *“Test suites should contain scalable functions”.*
 - (D) *“Test suites should contain problems with scalable evaluation cost”.*
 - (E) *“Test problems should have a canonical form”.*
- **BBOB** is an established testing framework for evaluating performance of continuous optimizers. The noise-free suite encompasses 24 functions.



The Traveling Salesman Problem

The *archetypical* Traveling Salesman Problem (TSP) is posed as finding a Hamilton circuit of minimal total cost. Explicitly, given a directed graph G , with a vertex set $V = \{1, \dots, |V|\}$ and an edge set $E = \{\langle i, j \rangle\}$, each edge has cost information $c_{ij} \in \mathbb{R}^+$.

Black-box formulation: permutations

$$\begin{array}{l}
 \text{[TSP-perm]} \quad \text{minimize} \quad \sum_{i=0}^{n-1} c_{\pi(i), \pi((i+1) \bmod n)} \\
 \text{subject to:} \\
 \pi \in P_{\pi}^{(n)}
 \end{array} \tag{1}$$

ILP Formulation [Miller-Tucker-Zemlin]

TSP as an ILP utilizes n^2 binary decision variables \mathbf{x}_{ij} :

$$\text{[TSP-ILP] minimize } \sum_{\langle i,j \rangle \in E} c_{ij} \cdot \mathbf{x}_{ij}$$

subject to:

$$\sum_{j \in V} \mathbf{x}_{ij} = 1 \quad \forall i \in V$$

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$$\mathbf{x}_{ij} \in \{0, 1\} \quad \forall i, j \in V$$

$$\mathbf{u}_i - \mathbf{u}_j + 1 \leq (|V| - 1)(1 - \mathbf{x}_{ij}) \quad \forall i, j \in 1 \dots |V|$$

$$|V| \geq \mathbf{u}_i \geq 2 \quad \forall i \in \{2, 3, \dots, |V|\}$$

(2)

where n integers \mathbf{u}_i are needed as decision variables to prevent inner-circles.

Q1: Problem Representation

[Q1] Should a **problem representation be dictated** per each benchmarking problem?

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[A1]

- A certain problem formulation **should be set fixed** — TSP-ILP and TSP-perm are two different search-problems!
- We suggest to restrict the benchmark suite to functions $f : \mathcal{S} \rightarrow \mathbb{R}^+$ (\mathcal{S} being a finite set of integers — also the most common representation in the EC literature)

Q2: Instance-Based Problems

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[A2]

- The hardness of an instance-based CO problem can differ substantially between two different instances
- Problem instances require specific descriptions and are therefore, in general, **not arbitrarily scalable with respect to their dimensionality**
- We suggest that **preference should be given to instance-free problems**; instance-based problems be included only to the extent needed to understand performance behavior that cannot be otherwise observed over instance-free problems

Q3: Invariant Problem Formulation

[Q3] Should the benchmarking framework cover the invariance aspect, and implicitly **favor algorithms that are invariant**?
If so, which invariances should be respected?

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[A3]

- Every benchmark suite might focus on a too narrow representative set of problems – with the risk of *overfitting*
- Therefore, the suite **should account for problem invariances**.
- We believe that conforming to certain, natural invariances reduces the risk of such overfitting. Our full paper elaborates on such meaningful invariances.

Q4-Q6: Performance Evaluation

[Q4] Which primary **performance evaluation measure** should be adopted?

[Q5] Should **performance aggregation** be conducted?

[Q6] Should the test-suite also facilitate **algorithm profiling** in the sense of algorithmic analysis beyond pure performance evaluation?

Performance Evaluation Answered

[A4] We advocate the use of **function evaluations as the main performance measure** — as in COCO.

[A5] Yes, we support **performance aggregation** (though not over problem dimensions, since it should be used for algorithm selection).

[A6] Yes, the benchmark suite **should allow for algorithm profiling**.

Admitting Runtime Analysis

- BBOB also encompasses a few rather simple problems like the Sphere function ($\mathbb{R}^n, F_1(x) := \sum_{i=1}^n x_i^2$) and other *unconstrained convex* problems.
- Convexity is irrelevant here, but the equivalent in problem hardness could be *simple problems admitting runtime analysis*.
- A well-known representative of this class is the “OneMax” problem

$$\begin{array}{l}
 \text{[HD] minimize } \sum_{i=1}^n x_i \\
 \text{subject to:} \\
 x_i \in \{0, 1\} \quad \forall i \in \{1, \dots, n\}
 \end{array}
 \tag{3}$$

Q7: Problems Admitting Runtime Analysis

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[A7]
Yes, selected analyzable functions, such as HD, **should be incorporated** into the test-suite, also to promote intensified discussions between theory-driven to practice-oriented scholars.

Q8: Facing Operations Research

Many CO problems may be formulated as Integer Linear Programs and treated by Mathematical Programming (MP) solvers in **extreme efficiency**.

Performance differences may be significant when compared to RSH.

[Q8] Should RSHs' performance be evaluated on **problems that are known to be effectively treated by MP-solvers** in practice?

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[Q8] Should RSHs' performance be evaluated on **problems that are known to be effectively treated by MP-solvers** in practice?

[A8]

Yes, problems that are (easily) solvable by MP-solvers could be incorporated into the test-suite, as long as they are not instance-based. Notably, a preference should be given to more challenging problems.

Constraints Satisfaction Problems

The n -queens problem (NQP) is defined as the task to place n queens on an $n \times n$ chessboard such that they cannot *capture* each other.

[NQP-CSP] satisfy:

$$\begin{aligned}
 & \sum_{i,j} x_{ij} = n \\
 & \sum_{i,j|j-i:=k} x_{ij} \leq 1 \quad k \in \{-n+2, -n+3, \dots, n-3, n-2\} \\
 & \sum_{i,j|i+j:=\ell} x_{ij} \leq 1 \quad \ell \in \{2, 3, \dots, 2n-3, 2n-2\} \\
 & x_{ij} \in \{0, 1\} \quad \forall i, j \in \{1, \dots, n\}
 \end{aligned} \tag{4}$$

n^2 binary decision variables x_{ij} are associated with the chessboard's coordinates, having an origin $(1, 1)$ at the top-left corner.

Q9: Distinguishing CSP?

The OR community distinguishes between standard optimization problems to Constraints Satisfaction Problems: Constraints Programming has forked into an independent subcommunity.

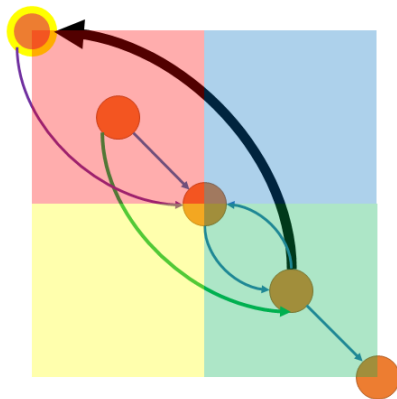
[Q9] Should RSHs' **performance be indistinguishably evaluated on CSPs as well**? That is, should a distinction between standard optimization to CSPs be avoided in the black-box perspective?

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[Q9] Should RSHs' **performance be indistinguishably evaluated on CSPs as well**? That is, should a distinction between standard optimization to CSPs be avoided in the black-box perspective?

[A9]
Yes, RSHs' performance **should be indistinguishably evaluated on CSPs as well**, since in the black-box perspective they are merely CO-problems.



Discussion

Nonlinear Hard Problems

This class of problems is meant to capture challenging CO problems that do not subscribe to MP/OR.

The Low-Autocorrelation Binary Sequence (LABS) problem is a hard CO problem with practical applications in electrical engineering. Given a sequence of length n , $S := (s_1, \dots, s_n)$ with $s_i = \pm 1$,

$$\begin{aligned}
 & \text{[LABS] maximize } \frac{n^2}{2E(S)} \\
 & \text{subject to:} \\
 & E(S) := \sum_{k=1}^{n-1} \left(\sum_{i=1}^{n-k} s_i \cdot s_{i+k} \right)^2 \\
 & s_i \in \{-1, +1\} \quad \forall i \in \{1 \dots n\}
 \end{aligned} \tag{5}$$

The Human Factor

The **human factor** plays a crucial role in such processes.

- Formulation of a test-suite may involve three types of scholars: theoreticians, algorithms' designers, and practitioners.
 - (i) theoreticians naturally favor analyzable functions
 - (ii) algorithms' engineers may prefer families of functions that are successfully treated by their designs
 - (iii) practitioners may have the best insights into which functions most accurately represent real-world problems (thus having their biased preferences)
- A proper balance should be made amongst those three parties to effectively compile a test-suite meaningful to a broad audience.

Communities and Resources

- **INFORMS**: The Institute for Operations Research and the Management Sciences; <https://www.informs.org/>
- **COIN-OR**: Computational Infrastructure for Operations Research – a project that aims to “create for mathematical software what the open literature is for mathematical theory”; <https://www.coin-or.org/>
- **MATHEURISTICS**: model-based metaheuristics, exploiting MP in a metaheuristic framework; <http://mh2018.sciencesconf.org/>

Benchmarking and Competitions

- MIPLIB: the Mixed Integer Programming LIBrary

<http://miplib.zib.de/>

- CSPLib: a problem library for constraints

<http://csplib.org/>

- SAT-LIB: the Satisfiability Library - Benchmark Problems

<http://www.cs.ubc.ca/~hoos/SATLIB/benchm.html>

- TSP-LIB: the Traveling Salesman Problem sample instances

<http://comopt.ifi.uni-heidelberg.de/software/TSPLIB95/>

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